

Number Theory

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• $\begin{cases} a \in \mathbb{Z} \\ b \in \mathbb{N} \end{cases}, \frac{a}{b} = \lfloor \frac{a}{b} \rfloor + d, (0 \leq d < 1) \Rightarrow a = b \lfloor \frac{a}{b} \rfloor + b d$
 $= b q + r, (0 \leq r < b)$

• $a = b q, b | a (b \neq 0)$

- 定理 (1) $a | b, b | a \Rightarrow a = \pm b$
 (2) $a | b, b | c \Rightarrow a | c$
 (3) $\begin{cases} d | a \\ d | b \end{cases} \Rightarrow$ (i) $\cancel{kd} | ka$ (ii) $d | ax + by$
 (ii) $d | a \pm b$

• 定義 (Module) $M \subset \mathbb{Z}, \forall a \in M \Rightarrow a \pm b \in M \quad (a \in M \Rightarrow \langle a \rangle \subset M)$

• 定理 M : module $\Rightarrow M = \langle d \rangle, \quad \text{証 (i) } M = \{0\} \text{ 或 } 0 \leq r < d$
 (ii) $d = \min_{a \in M, a > 0} a, \forall n \in M, n = dq + r \Rightarrow r = 0$

- 例 (1) $\langle a \rangle = a\mathbb{Z} = \{ax | x \in \mathbb{Z}\} = \{\dots, -2a, -a, 0, a, 2a, \dots\}$
 (2) $\langle a, b \rangle = a\mathbb{Z} + b\mathbb{Z} = \{ax + by | x, y \in \mathbb{Z}\} = \langle d \rangle$
 (3) $\langle a_1, \dots, a_k \rangle = a_1\mathbb{Z} + \dots + a_k\mathbb{Z} = \{a_1x_1 + \dots + a_kx_k | x_i \in \mathbb{Z}\}$

(最大公因子)

- 定理 (1) $\exists x, y, d = ax + by$
 (2) $d | a, d | b \quad \therefore d = (a, b)$
 (3) $e | a, e | b \Rightarrow e | d$

• Euclid algorithm $(a, b) = (b, r), a = bq + r$

$a = 323, b = 221$
 $323 = 1 \cdot 221 + 102$
 $221 = 2 \cdot 102 + 17$
 $102 = 6 \cdot 17$

1	323	221	2
	221	204	
6	102	17	
	102		
	0		

• 定理 $ax + by = n$ 有整數解 $\Leftrightarrow (a, b) | n$

質數, 互質 ($a \perp b$)

$\Rightarrow 17 = b - 2 \cdot 102$
 $= b - 2(a - b)$
 $= -2a + 3b$

• $(323, 221)$
 $= (221, 102)$
 $= (102, 17) = 17$
 • $(-2) \cdot 323 + 3 \cdot 221 = 17$

• 定理 (Euclid) ∞ 個 primes

• 定理 $a^m - 1$ prime $\Rightarrow \begin{cases} a = 2 \\ m = \text{prime} \end{cases}$ (Mersenne)

• 定理 $2^m + 1$ " $\Rightarrow m = 2^k$ (Fermat)

• 定理 (Prime th) $\pi(x) \stackrel{\text{def}}{=} (\# \text{ primes } \leq x) \approx \frac{x}{\ln x}$

(i) $\lim_{x \rightarrow \infty} \pi(x) = \infty$

(ii) $\lim_{x \rightarrow \infty} \frac{\pi(x)}{x} = 0$

(iii) n^{th} prime $P_n \approx n \ln n$

• 解 $323x + 221y = 85 \quad 85 = 17 \cdot 5$

(甲) $323(-10) + 221 \cdot 15 = 85$ (特別解)

(乙) $19x + 13y = 5 \quad (19, 13) = 1$

$19(-2) + 13 \cdot 3 = 1$

$19(-10) + 13(15) = 5$ (一般解)

$\underline{+13t} \quad \underline{-19t}$

• 定理 (1) $d | ab, (d, a) = 1 \Rightarrow d | b$

定理 (算數基本定理)

(2) prime $p | ab \Rightarrow p | a$ 或 $p | b \quad n = p_1^{r_1} \dots p_r^{r_r}$ (唯一分解)

証: (1) $dx + ay = 1 \Rightarrow dbx + aby = b$

同餘 (Congruence)

$a \equiv b \pmod{m} \iff m \mid a-b$

- 定理 (1) $a \equiv a \pmod{m}$
 - (2) $a \equiv b \pmod{m} \Rightarrow b \equiv a \pmod{m}$
 - (3) $\begin{cases} a \equiv b \pmod{m} \\ b \equiv c \pmod{m} \end{cases} \Rightarrow a \equiv c \pmod{m}$
- Equivalence class $= \{[0], [1], \dots, [m-1]\} = \mathbb{Z}_m$
 $[a] = \{k \equiv a \mid k \in \mathbb{Z}\} = a + m\mathbb{Z}$

定理 (1) $\begin{cases} a \equiv c \pmod{m} \\ b \equiv d \pmod{m} \end{cases} \Rightarrow \begin{cases} a \pm b \equiv c \pm d \\ ka \equiv kc \pmod{m} \\ ab \equiv cd \end{cases}$ 証 (1) $\begin{cases} m \mid a-c \\ m \mid b-d \end{cases} \Rightarrow m \mid (a+b)-(c+d)$
 $ab \equiv cb \equiv cd \pmod{m}$

(2) $\begin{cases} (i) ka \equiv kb \pmod{m}, k \perp m \Rightarrow a \equiv b \pmod{m} \\ (ii) ka \equiv kb \pmod{m} \begin{cases} k = k_0 d \\ m = m_0 d \\ k_0 \perp m_0 \end{cases} \Rightarrow a \equiv b \pmod{m_0} \end{cases}$ (2) $m \mid k(a-b)$
 $m_0 \mid k_0 d (a-b)$

例 $\begin{cases} 13 \cdot 5 \equiv 7 \cdot 5 \pmod{6} \Rightarrow 13 \equiv 7 \pmod{6} \\ 5 \cdot 10 \equiv 14 \cdot 10 \pmod{6} \Rightarrow 5 \equiv 14 \pmod{3} \end{cases}$ (Euler) $\varphi(m) := \# a, \begin{matrix} 1 \leq a < m \\ a \perp m \end{matrix}$

- 定理 $\begin{cases} \mathbb{Z}_m = \{[0], [1], \dots, [m-1]\} & |\mathbb{Z}_m| = m \\ \mathbb{Z}_m^* = \{[a] \mid 0 \leq a \leq m-1, a \perp m\} & |\mathbb{Z}_m^*| = \varphi(m) \end{cases}$
- (1) $(\mathbb{Z}_m, +, \cdot)$ 環, (2) (\mathbb{Z}_m^*, \cdot) 乘法群, (3) \mathbb{Z}_p 有限域

$\mathbb{Z}_{12}: \begin{cases} 9 \equiv 33 \pmod{12} \\ 5 \equiv 17 \pmod{12} \end{cases} \begin{cases} 9+5 \equiv 33+17 \equiv 2 \\ 9-5 \equiv 33-17 \equiv 4 \pmod{12} \\ 9 \cdot 5 \equiv 33 \cdot 17 \equiv 9 \end{cases}$
 $\begin{cases} [a] + [0] = [a] \\ [a] \cdot [1] = [a] \end{cases} \quad [9] + [5] = [2], [9] - [5] = [4], [9][5] = [9]$

- 定義 $\{x_1, \dots, x_m\}$ ($x_i \not\equiv x_j \pmod{m}$), Complete residue system (= \mathbb{Z}_m)
- $\{x_1, \dots, x_{\varphi(m)}\}$ ($\forall x_i \perp m$), reduced " " (= \mathbb{Z}_m^*)

定理 $a \perp m \begin{cases} (1) \{x_1, \dots, x_m\} \text{ complete r.s.} \Rightarrow \{ax_1, \dots, ax_m\} \text{ also, } a\mathbb{Z}_m \equiv \mathbb{Z}_m \pmod{m} \\ (2) \{x_1, \dots, x_{\varphi(m)}\} \text{ reduced " "} \Rightarrow \{ax_1, \dots, ax_{\varphi(m)}\} \text{ " , } a\mathbb{Z}_m^* \equiv \mathbb{Z}_m^* \end{cases}$

証 (1) $ax_i \equiv ax_j \pmod{m} \Rightarrow m \mid a(x_i - x_j) \Rightarrow m \mid x_i - x_j \Rightarrow x_i \equiv x_j$ (2) $ax_i \perp m$
 (1) $(a, m) = d, ax \equiv ay \pmod{m} \Rightarrow x \equiv y \pmod{m_0}$

例 $5\mathbb{Z}_{12} = \{0, 5, 10, 15, 20, 25, 30, 35, 40, 45, 50, 55\} \equiv \{0, 5, 10, 3, 8, 1, 6, 11, 4, 9, 2, 7\} \pmod{12}$
 $5\mathbb{Z}_{12}^* = \{5, 25, 35, 55\} \pmod{12} \equiv \{5, 1, 11, 7\}$

$10\mathbb{Z}_{12} = \{0, 10, 20, 30, 40, 50, 60, 70, 80, 90, 100, 110\} \equiv \{0, 10, 8, 6, 4, 2, 0, 10, 8, 6, 4, 2\} \pmod{12} = 2 * \{0, 5, 4, 3, 2, 1\} \Rightarrow a \equiv b \pmod{6}$
 $(10, 12) = 2$

• 定理 (Euler) $a \perp m, a^{\varphi(m)} \equiv 1 \pmod{m}$ 証: $x_1 \dots x_{\varphi(m)} \equiv a^{\varphi(m)} x_1 \dots x_{\varphi(m)} \pmod{m}$

• 定理 (Fermat) $p: \text{prime}, a^p \equiv a \pmod{p}$ $\begin{cases} \text{(i) } p|a, a \equiv a^p \equiv 0 \pmod{p} \\ \text{(ii) } p \nmid a, a^{p-1} \equiv 1 \pmod{p} \end{cases}$

• 定理 (Wilson) $p: \text{prime} \iff (p-1)! \equiv -1 \pmod{p}$ $p=11$
証 " \Rightarrow " $\forall 1 \leq k \leq p-1 \begin{cases} k x_i \equiv 1 \pmod{p} \\ k k \equiv 1 \pmod{p} \end{cases} \Rightarrow k = \pm 1$ $\{1, 2, \overbrace{3, 4}, \overbrace{5, 6}, \overbrace{7, 8}, 9, 10\}$

" \Leftarrow " $d|p \Rightarrow (p-1)! \equiv -1 \pmod{d} \Rightarrow 0 \equiv -1 \pmod{d} \Rightarrow d=1 \neq p$
 $(1 \leq d < p)$

• 定義 $g: \text{積性函數}$: $m \perp n \Rightarrow g(mn) = g(m)g(n) \iff \begin{cases} g(p_1^{k_1} \dots p_r^{k_r}) = g(p_1^{k_1}) \dots g(p_r^{k_r}) \\ g(1) = 1 \end{cases}$
(multiplicative)

• (Euler-totient) $\varphi(n) := \#d, 1 \leq d \leq n, d \perp n$

(1) φ : 積性

(2) $p: \text{prime}, \varphi(p^k) = p^k - p^{k-1} = p^k(1 - \frac{1}{p})$

(3) $m = p_1^{k_1} \dots p_r^{k_r}, \varphi(m) = m(1 - \frac{1}{p_1}) \dots (1 - \frac{1}{p_r})$

(4) $\sum_{d|m} \varphi(d) = m$

(1)+(2) \iff (3)

(*) (甲) \iff (乙) \iff (丙) \iff (丁)

(*) 定理 2 + 推論, $p=5$

証 (3) $m = 2^3 \cdot 3^2 \cdot 5, \varphi(m) = m - \frac{m}{2} - \frac{m}{3} - \frac{m}{5} + \frac{m}{2 \cdot 3} + \frac{m}{3 \cdot 5} + \frac{m}{2 \cdot 5} - \frac{m}{2 \cdot 3 \cdot 5}$

$$= m(1 - \frac{1}{2})(1 - \frac{1}{3})(1 - \frac{1}{5})$$

$$d \in (1+2+2^2+2^3)(1+3+3^2)(1+5)$$

(4) " (甲) $\sum_{d|m} \varphi(d) = (1 + \varphi(2) + \varphi(2^2) + \varphi(2^3))(1 + \varphi(3) + \varphi(3^2))(1 + \varphi(5))$

$$= (1 + (2-1) + (2^2-2) + (2^3-2^2))(1 + (3-1) + (3^2-3))(1 + (5-1))$$

$$= 2^3 \cdot 3^2 \cdot 5$$

(乙) $\{ \frac{1}{12}, \frac{1}{12}, \frac{2}{12}, \frac{3}{12}, \frac{4}{12}, \frac{5}{12}, \frac{6}{12}, \frac{7}{12}, \frac{8}{12}, \frac{9}{12}, \frac{10}{12}, \frac{11}{12} \}$ $m=12$

(約分) $\Rightarrow \{ \frac{1}{1}, (\frac{1}{2}), (\frac{1}{3}, \frac{2}{3}), (\frac{1}{4}, \frac{3}{4}), (\frac{1}{6}, \frac{5}{6}), (\frac{1}{12}, \frac{5}{12}, \frac{7}{12}, \frac{11}{12}) \}$

$$\varphi(1) + \varphi(2) + \varphi(3) + \varphi(4) + \varphi(6) + \varphi(12) = 12$$

• (Möbius) $\mu(n) = \mu(p_1^{k_1} \dots p_r^{k_r}) := \begin{cases} (-1)^r & k_1 = k_2 = \dots = k_r = 1 \\ 0 & \text{else} \end{cases} \quad (\mu(1) = 1)$

(1) μ : 積性

証: (1) $\mu(p_1^{k_1} \dots p_r^{k_r} \cdot g_1^{h_1} \dots g_s^{h_s}) = \mu(p_1^{k_1} \dots p_r^{k_r}) \mu(g_1^{h_1} \dots g_s^{h_s})$

(2) $\sum_{d|m} \mu(d) = \delta_{m=1}$

(2) $\sum_{d|m} \mu(d) = 1 + \mu(2) + \mu(3) + \mu(5) + \mu(2 \cdot 3) + \mu(3 \cdot 5) + \mu(2 \cdot 5) + \mu(2 \cdot 3 \cdot 5)$

$$m = 2^3 \cdot 3^2 \cdot 5 = 1 + \binom{3}{1}(-1) + \binom{3}{2}(-1)^2 + \binom{3}{3}(-1)^3 = (1-1)^3 = 0$$

Möbius Transformation

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- 定義 (1) g 積性函數: $m \perp n, g(mn) = g(m)g(n) \Leftrightarrow \begin{cases} g(1) = 1 \\ g(p_1^{k_1} \dots p_r^{k_r}) = g(p_1^{k_1}) \dots g(p_r^{k_r}) \end{cases}$
- (2) Möbius 函數 μ : $\sum_{d|m} \mu(d) = \delta_{m=1} = \begin{cases} 1 & m=1 \\ 0 & \text{else} \end{cases}$ (recursive 定義)



• Möbius transf: $f \longrightarrow g(m) = \sum_{d|m} f(d)$

m	1	2	3	4	5	6	7	8	9	10	11	12	15	30
$\mu(m)$	1	-1	-1	0	-1	1	-1	0	0	1	-1	0	1	-1

$\mu \rightarrow \delta_{m=1}$

$\varphi \rightarrow m$

$1 \rightarrow \tau(m) = \# \text{ 因數}, = 4 \cdot 3 \cdot 2$

$m \rightarrow \sigma(m) = \text{因數和}, = (1+2+2^2+2^3)(1+3+3^2)(1+5) \quad m = 2^3 \cdot 3^2 \cdot 5$

$m = p^k, 0 = 1 + \mu(p) + \mu(p^2) + \dots + \mu(p^k)$

$\mu(p) = -1, \mu(p^2) = \mu(p^3) = \dots = 0$

(Möbius Inversion)

• 定理 1 $g(m) = \sum_{d|m} f(d) \Leftrightarrow f(m) = \sum_{d|m} \mu(d) g(\frac{m}{d}) \left[= \sum_{d|m} \mu(d) g(k) = \sum_{d|m} \mu(\frac{m}{d}) g(d) \right]$

• 定理 2 f 積性 $\Leftrightarrow g$ 積性

• 推論

(1) $\delta_{m=1}, m, 1$ 積性 $\Rightarrow \mu, \varphi, \tau, \sigma$ 積性

(2) $\delta_{m=1} = \sum_{d|m} \mu(d) \Rightarrow \mu(m) = \begin{cases} (-1)^r & m = p_1 \dots p_r \\ 0 & \text{else} \end{cases} \quad \begin{aligned} \mu(m) &= \mu(p_1^{k_1} \dots p_r^{k_r}) \\ &= \mu(p_1^{k_1}) \dots \mu(p_r^{k_r}) \end{aligned}$

$(m = 2^3 \cdot 3^2 \cdot 5^1)$

(3) $m = \sum_{d|m} \varphi(d) \Rightarrow \varphi(m) = \sum_{d|m} \mu(d) \frac{m}{d} = m \left[1 - \frac{1}{2} - \frac{1}{3} - \frac{1}{5} + \frac{1}{2 \cdot 3} + \frac{1}{2 \cdot 5} + \frac{1}{3 \cdot 5} - \frac{1}{2 \cdot 3 \cdot 5} \right]$
 $= m \left(1 - \frac{1}{p_1} \right) \dots \left(1 - \frac{1}{p_r} \right) = m \left(1 - \frac{1}{2} \right) \left(1 - \frac{1}{3} \right) \left(1 - \frac{1}{5} \right)$

(4) 環狀項鍊 $\begin{cases} m \text{ beads} \\ c \text{ colors} \end{cases}$ 個數 $N(m, c) = ?$

$f(x+d) = f(x), \text{ period} = d$

$N(d) = \# \text{ with period } d$

例 $m=4, C=2 (0,1)$

• $c^m = \sum_{d|m} d N(d)$ (直線排列)

• $2^4 = 16$

• $m N(m) = \sum_{d|m} \mu(d) c^{\frac{m}{d}}$

• $d=1, N(1) = 2 \quad \begin{cases} 0000 \\ 1111 \end{cases}$

$d=2, N(2) = 1 \quad \begin{cases} 0101 \rightarrow 1010 \end{cases}$

• $N(m, c) = \sum_{k|m} N(k)$

$d=4, N(4) = 3 \quad \begin{cases} 0111 \rightarrow 1011, 1101, 1110 \\ 0011 \rightarrow 1001, 1100, 0110 \\ 0001 \rightarrow 1000, 0100, 0010 \end{cases}$

$= \sum_{k|m} \frac{1}{k} \sum_{d|k} \mu(d) c^{\frac{k}{d}}$

$= \frac{1}{m} \sum_{d|m} \varphi(d) c^{\frac{m}{d}}$

• $2^4 = 1 \cdot N(1) + 2 \cdot N(2) + 4 \cdot N(4)$

• $N(1) = \mu(1) 2^1 = 2$

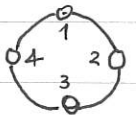
• $N(2) = \frac{1}{2} \left[\mu(1) 2^2 + \mu(2) 2^1 \right] = 1$

• $N(4) = \frac{1}{4} \left[\mu(1) 2^4 + \mu(2) 2^2 + \mu(4) 2^1 \right] = 3$

• $N(4, 2) = \frac{1}{4} \left[\varphi(1) 2^4 + \varphi(2) 2^2 + \varphi(4) 2^1 \right]$

$= 6$

• $N(4, 2) = N(1) + N(2) + N(4) = 6$



定理証明

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• 定理 1 $g(m) = \sum_{d|m} f(d) \Leftrightarrow f(m) = \sum_{d|m} \mu(d) g\left(\frac{m}{d}\right) \quad \left(\sum_{d|m} \mu(d) = \delta_{m=1}\right)$

証 " \Rightarrow " $\sum_{d|12} \mu(d) g\left(\frac{12}{d}\right) \stackrel{m=12}{=} \mu(1)g(12) + \mu(2)g(6) + \mu(3)g(4) + \mu(4)g(3) + \mu(6)g(2) + \mu(12)g(1)$

$$= \mu(1) [f(1) + f(2) + f(3) + f(4) + f(6) + f(12)]$$

$$+ \mu(2) [f(1) + f(2) + f(3) + f(6)]$$

$$+ \mu(3) [f(1) + f(2) + f(4)]$$

$$+ \mu(4) [f(1) + f(3)]$$

$$+ \mu(6) [f(1) + f(2)]$$

$$+ \mu(12) [f(1)]$$

$$= f(12)$$

" \Leftarrow " $\sum_{d|12} f(d) \stackrel{m=12}{=} f(1) + f(2) + f(3) + f(4) + f(6) + f(12)$

$$= \left[\begin{array}{l} \mu(1)g(1) \\ \mu(1)g(2) + \mu(2)g(1) \\ \mu(1)g(3) + \mu(3)g(1) \\ \mu(1)g(4) + \mu(2)g(2) + \mu(4)g(1) \\ \mu(1)g(6) + \mu(2)g(3) + \mu(3)g(2) + \mu(6)g(1) \\ \mu(1)g(12) + \mu(2)g(6) + \mu(3)g(4) + \mu(4)g(3) + \mu(6)g(2) + \mu(12)g(1) \end{array} \right]$$

$$= g(12)$$

• 定理 2 $f: \text{積性} \Leftrightarrow g \text{積性} \quad \left\{ \begin{array}{l} \forall ab < n, a \perp b \\ f(ab) = f(a)f(b) \end{array} \right.$

証 " \Leftarrow " 已知 g 積性, 擬証 f 積性 for $n (=12)$ by induction: $g(2^2 \cdot 3) = g(2^2)g(3) \stackrel{?}{\Rightarrow} f(2^2 \cdot 3) = f(2^2)f(3)$

$$g(2^2 \cdot 3) = \sum_{d|12} f(d) = \sum_{\substack{a|2^2 \\ b|3}} f(ab) = f(1 \cdot 1) + f(1 \cdot 3) + f(2 \cdot 1) + f(2 \cdot 3) + f(2^2 \cdot 1) + \underline{f(2^2 \cdot 3)}$$

$$g(2^2)g(3) = \left(\sum_{a|2^2} f(a)\right) \left(\sum_{b|3} f(b)\right) = [f(1) + f(2) + f(2^2)] [f(1) + f(3)]$$

$$= f(1)f(1) + f(1)f(3) + f(2)f(1) + f(2)f(3) + f(2^2)f(1) + \underline{f(2^2)f(3)}$$

$$\Rightarrow f(2^2 \cdot 3) = f(2^2)f(3)$$

" \Rightarrow "

$$(d|2^2 \cdot 3 : (1+2+2^2) \cdot (1+3))$$

$$d = a \cdot b$$

Stern-Brocot tree : Binary search tree for reduced fractions (最約分數)

(1) 左右反对稱

(2) Binary search tree:

$$\frac{m}{n} < \frac{m+m'}{n+n'} < \frac{m'}{n'} \quad (mn' < nm')$$

$$(3) \begin{vmatrix} n & n' \\ m & m' \end{vmatrix} = 1$$

$$\Rightarrow \begin{vmatrix} n & n+n' \\ m & m+m' \end{vmatrix} = \begin{vmatrix} n+n' & n' \\ m+m' & m' \end{vmatrix} = 1$$

(4) $n \perp m$

(5) $\frac{m}{n} \in \mathbb{Q}^+$, 恰一次 (i) 唯一 (ii) 存在.

(6) $\mathbb{Q}^+ \longleftrightarrow \{L, R\}^*$ (number system)

$$" \leftarrow " \quad LRRL = \begin{bmatrix} 1 & 1 \\ 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 \\ 1 & 1 \end{bmatrix}^2 \begin{bmatrix} 1 & 1 \\ 0 & 1 \end{bmatrix} = \begin{bmatrix} 3 & 4 \\ 2 & 3 \end{bmatrix}$$

$$f: \begin{bmatrix} 3 & 4 \\ 2 & 3 \end{bmatrix} \rightarrow \frac{5}{7}$$

" \rightarrow " $\frac{m}{n} \rightarrow ?$

(甲) Binary search:

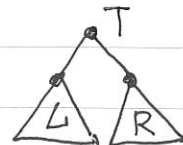
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• S := I;
  while m/n ≠ f(S) do
    if m/n < f(S) then (output(L); S := SL)
    else (output(R); S := SR)
  
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(乙) subtrees 相似性

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• while m ≠ n do
  if m < n then (output(L); n := n - m)
  else (output(R); m := m - n)
  
```



• $m/n = 5/7$

• $m = 5 \quad 5 \quad 3 \quad 1 \quad 1$

$n = 7 \quad 2 \quad 2 \quad 2 \quad 1$

output L R R L

$$L \leftrightarrow T \leftrightarrow R$$

$$\frac{m}{n+m} \quad \frac{m}{n} \quad \frac{n+m}{n}$$

$$RRL = \frac{5}{7} \quad \frac{5}{2} \quad \frac{7}{2}$$

